

Higher-order Learning in Multi-Agent Games

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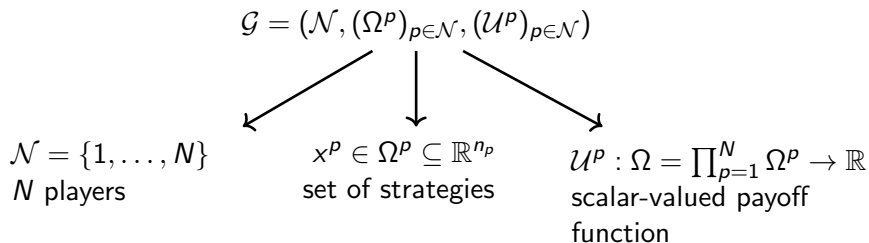
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*based on joint work with Bolin Gao

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Game Setup



Assumption. \mathcal{G} is **continuous concave**, i.e., $\forall p$,

- Ω^p non-empty, closed, convex subset of \mathbb{R}^{n_p} ,
- $\mathcal{U}^p(x^p; x^{-p}) = \mathcal{U}^p(x)$ (jointly) continuous in $x = (x^p; x^{-p})$,
- $\mathcal{U}^p(x^p; x^{-p})$ concave and \mathcal{C}^1 in x^p , $\forall x^{-p} \in \Omega^{-p}$.

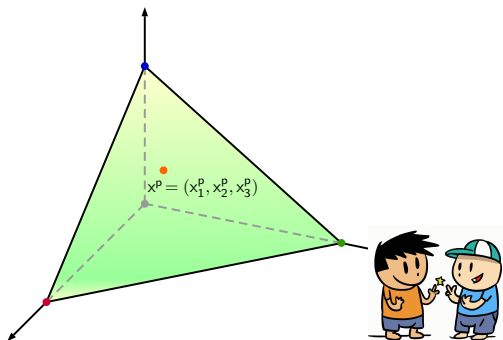
Example 1: Rock-Paper-Scissors

$$\begin{array}{c} \mathbf{r} \\ \mathbf{p} \\ \mathbf{s} \end{array} \begin{bmatrix} \mathbf{r} & \mathbf{p} & \mathbf{s} \\ (0, 0) & (-l, w) & (w, -l) \\ (w, -l) & (0, 0) & (-l, w) \\ (-l, w) & (w, -l) & (0, 0) \end{bmatrix}$$

l : loss > 0 , w : win > 0

- $i^P \in \{\mathbf{r}, \mathbf{p}, \mathbf{s}\}$
- $\mathbf{i} = (i^1; i^2) \sim$ reward
- x^P prob. of selection
- \mathcal{U}^P expected reward,

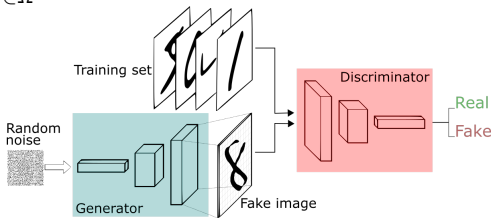
$$\mathcal{U}^P(x) = \pm x^{P\top} \begin{bmatrix} 0 & -l & w \\ w & 0 & -l \\ -l & w & 0 \end{bmatrix} x^{-P}$$



Example 2: Saddle-point problems

$$\max_{x^1 \in \Omega^1} \min_{x^2 \in \Omega^2} f(x^1, x^2)$$

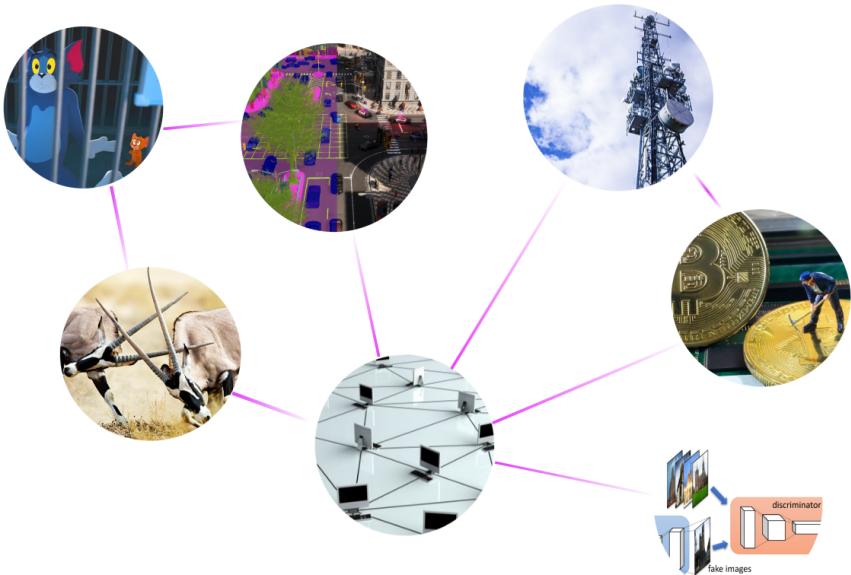
- designer and critic
- designer: submits design, x^1
- critic: submits appraisal, x^2
- $U^1 = f(x^1, x^2) = -U^2$
"agreement/satisfaction"



Generative Adversarial Network (GAN)

After a "few" simplifying assumptions:

$$f(x^1, x^2) = x^1 \cdot x^2$$



Behavioral science, ecology, wireless networks, (virtual) economy, traffic modeling...

Game Solution

Player p 's goal: given $x^{-p} \in \Omega^{-p}, \forall p \in \mathcal{N}$, $\max_{x^p \in \Omega^p} \mathcal{U}^p(x^p; x^{-p})$

$x^* = (x^{p*}; x^{-p*}) \in \Omega$ is **Nash equilibrium** (NE) when no players can benefit from unilateral deviation:

$$\mathcal{U}^p(x^{p*}; x^{-p*}) \geq \mathcal{U}^p(x^p; x^{-p*}), \forall x^p \in \Omega^p, \forall p \in \mathcal{N} \quad (1)$$

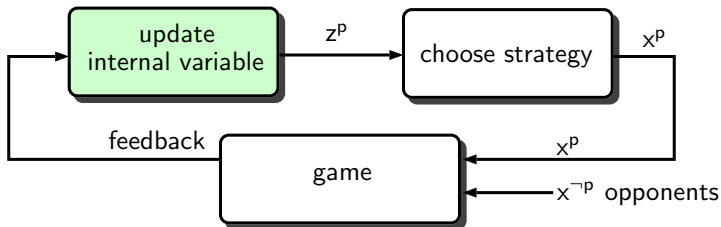
Equiv., under our concave game assumption,

$$(x - x^*)^\top U(x^*) \leq 0, \forall x \in \Omega \quad (2)$$

$U(x) = (U^p)_{p \in \mathcal{N}} = (\nabla_{x^p} \mathcal{U}^p(x^p; x^{-p}))_{p \in \mathcal{N}}$ (**pseudo-gradient**).

Players use simple rules/models to convert game information to their strategies, hopefully leading to a NE.

NE seeking via online learning

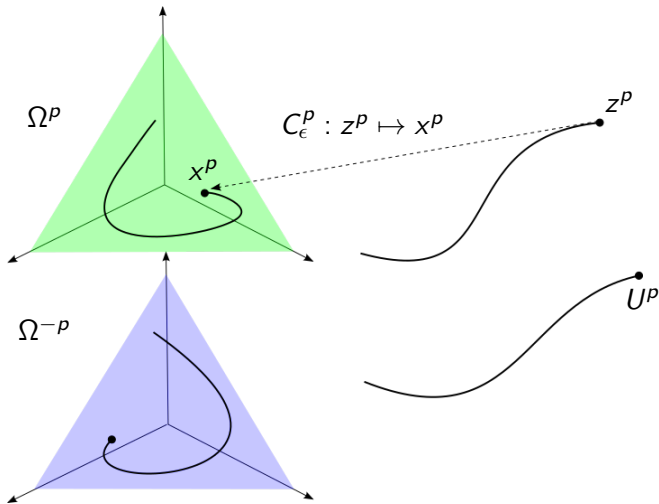


A group of players **learns** a NE by individually:

- choose strategies via an internal variable[†]
- receive information (feedback) from the game
- update own internal variable (\leftarrow “learning”)

[†]Q-value, dual aggregate, score, perception, model weights, etc.

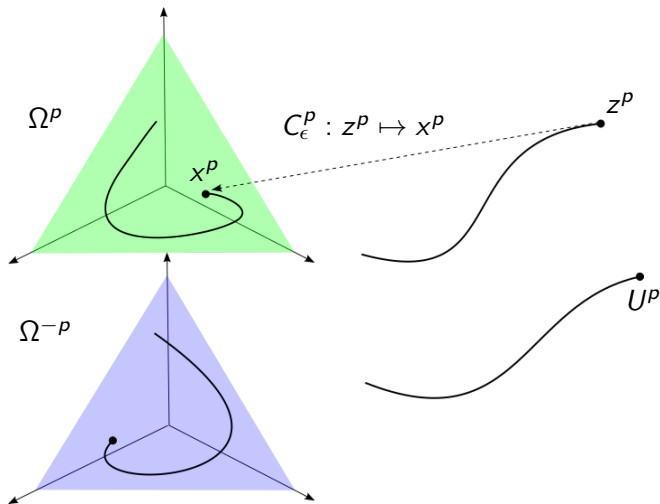
Each player maps its own variable z^P into a strategy $x^P \in \Omega^P$ through a *mirror map* $C_\epsilon^P : \mathbb{R}^{n_P} \rightarrow \Omega^P$



$$C_\epsilon^P(z^P) = \operatorname{argmax}_{y^P \in \Omega^P} \left[y^{P\top} z^P - \epsilon v^P(y^P) \right], \epsilon > 0 \quad (*)$$

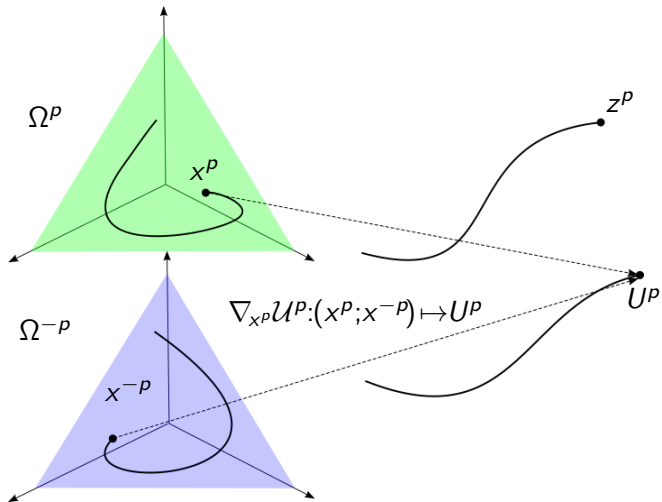
v^P : regularizer, ϵ : regularization constant

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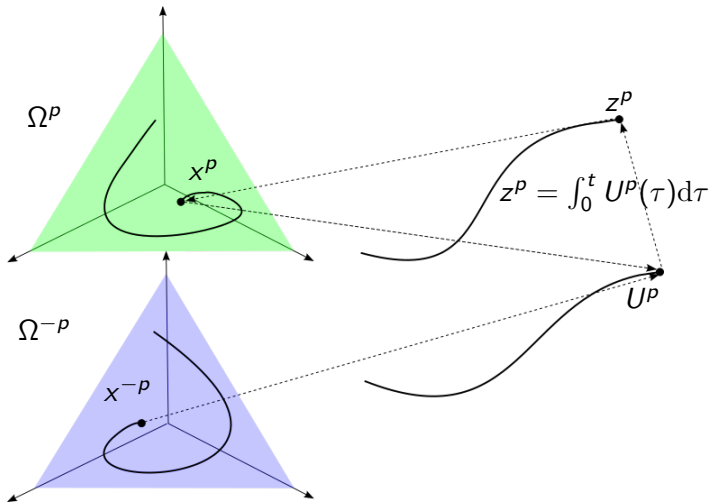


- $v^P = \|x^P\|_2^2/2 \Rightarrow C_\epsilon^P(z^P) = \pi_{\Omega^P}(\epsilon^{-1}z^P)$ (Projection)
- $v^P = x^P{}^\top \log(x^P) \Rightarrow C_\epsilon^P(z^P) = \frac{\exp(\epsilon^{-1}z^P)}{\sum_{q \in \mathcal{N}} \exp(\epsilon^{-1}z^q)}$ (Softmax)

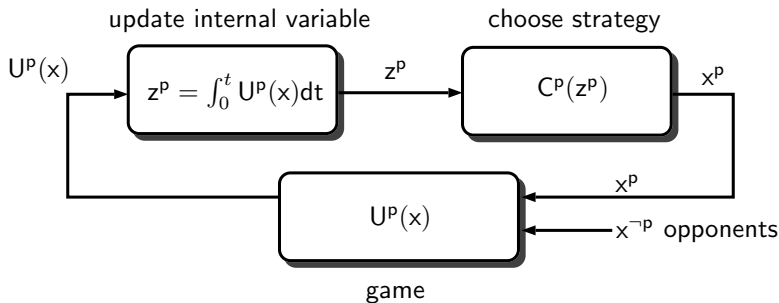
Each player plays the game \mathcal{G} using strategy x^P (taking into account opponents' $x^{-P} \in \Omega^{-P}$) and obtains partial-gradient U^P



Each player maps U^P back into z^P via aggregation
 $z^P = \int_0^t U^P(\tau) d\tau$ and the process continues indefinitely



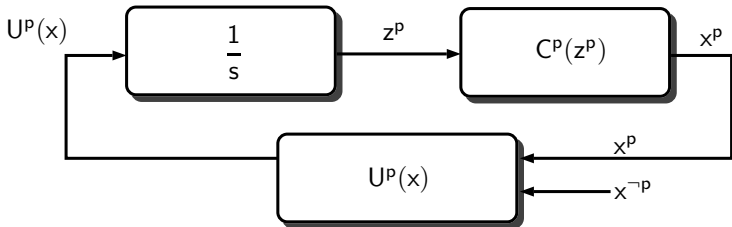
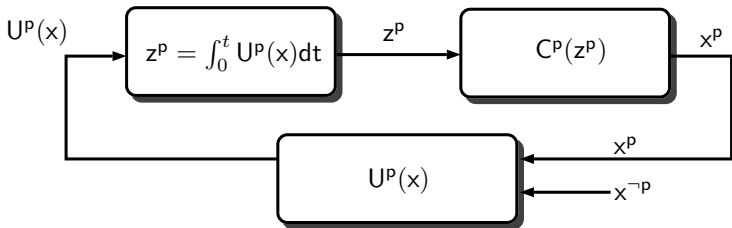
(Continuous-Time) Mirror Descent (Nemirovski '83)

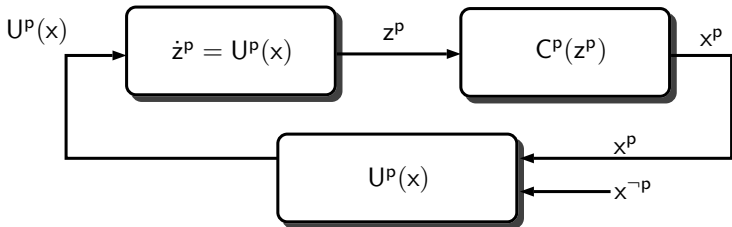
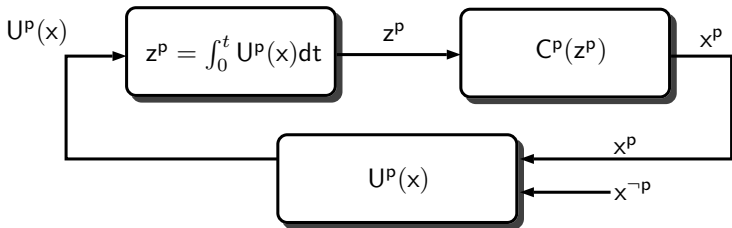


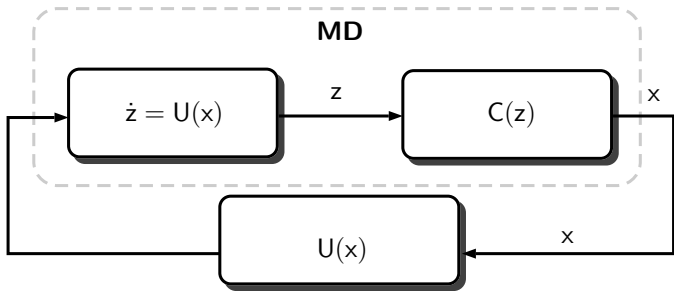
- z^p : a vector of player p 's internal state
- C^p : "mirror map", $z^p \mapsto x^p$ e.g., Euclidean projection
- U^p : a partial-gradient

$$U^p(x) = \nabla_{x^p} \mathcal{U}^p(x^p; x^{-p}) = \frac{\partial \mathcal{U}^p(x^p; x^{-p})}{\partial x^p}$$

- $x = (x^p; x^{-p})$: joint strategy







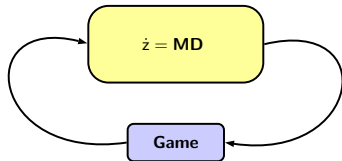
To analyze collective behavior,
stack all $p = 1, \dots, N$:

$$z = \begin{bmatrix} z^1 \\ \vdots \\ z^N \end{bmatrix} \quad U = \begin{bmatrix} U^1 \\ \vdots \\ U^N \end{bmatrix}$$

U is called a **pseudo-gradient**.

$$x = \begin{bmatrix} x^1 \\ \vdots \\ x^N \end{bmatrix} \quad C = \begin{bmatrix} C^1 \\ \vdots \\ C^N \end{bmatrix}$$

Convergence of MD



Using stacked notation,

$$z = \begin{bmatrix} z^1 \\ \vdots \\ z^N \end{bmatrix} \quad U = \begin{bmatrix} U^1 \\ \vdots \\ U^N \end{bmatrix} \quad x = \begin{bmatrix} x^1 \\ \vdots \\ x^N \end{bmatrix} \quad C_\epsilon = \begin{bmatrix} C_\epsilon^1 \\ \vdots \\ C_\epsilon^N \end{bmatrix}$$

Can represent this entire process as,

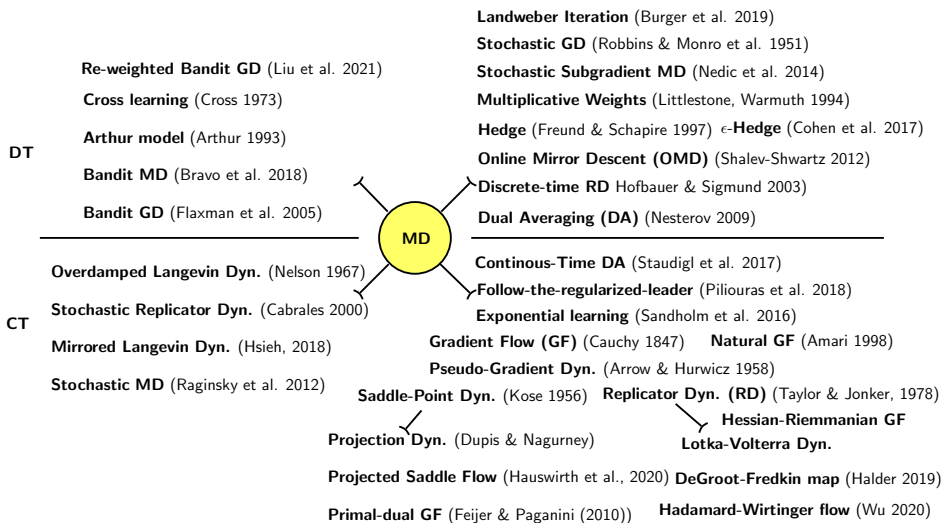
$$z = \int_0^t U(x) d\tau, \quad x = C_\epsilon(z), \quad (3)$$



$$\dot{z} = U(x), \quad x = C_\epsilon(z),$$

MD

Question: when does $x(t) = C_\epsilon(z(t))$ converge to x^* ?



MD converges under one of a trio of strict assumptions ¹ 8/

1. \mathcal{G} is **strictly monotone** (Rosen '65),

$$(U(x) - U(x'))^\top (x - x') < 0, \forall x \in \Omega \setminus \{x'\} \quad (4)$$

2. x^* is a **strict variationally stable state (VSS)** (Smith '73),

$$U(x)^\top (x - x^*) < 0, \forall x \in \Omega \setminus \{x^*\} \quad (5)$$

NE of strictly monotone game \implies strict VSS

3. x^* is a **strict Nash equilibrium** (Harsanyi '73),

$$U^p(x^p; x^{-p*}) < U^p(x^{p*}; x^{-p*}), \forall x^p \in \Omega^p \setminus \{x^{p*}\}, \forall p \in \mathcal{N} \quad (6)$$

or,

$$U(x^*)^\top (x - x^*) < 0, \forall x \in \Omega \setminus \{x^*\} \quad (7)$$

strict NE = strict VSS in finite games

¹See e.g., Mertikopoulos & Sandholm '16, Mertikopoulos & Staudigl '17, Mertikopoulos & Zhou '19, Migot & Cojocaru '21, Laraki & Mertikopoulos '13, Higher order game dynamics

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Two wider settings MD can fail to converge

1. \mathcal{G} is **merely monotone** (= stable game) (Hofbauer et al. '09)

$$(U(x) - U(x'))^\top (x - x') \leq 0, \forall x \in \Omega \quad (8)$$

2. x^* is a **mere variationally stable state** (VSS) (Smith, '82)

$$U(x)^\top (x - x^*) \leq 0, \forall x \in \Omega \quad (9)$$

NE of merely monotone game \implies mere VSS

Example. (Two-Player Rock-Paper-Scissors (RPS))

$$\mathcal{A} = \begin{bmatrix} 0 & -1 & 1 \\ 1 & 0 & -1 \\ -1 & 1 & 0 \end{bmatrix} \quad \mathcal{U}^p(x^p; x^q) = x^{p\top} \mathcal{A} x^q, x^p \in \mathbf{Simplex}$$

RPS with \mathcal{A} is merely monotone \implies unique interior NE

$x^* = (x^{p*}), x^{p*} = (1/3, 1/3, 1/3)$ is a mere VSS.

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Example. (Non-Monotone Game with Mere VSS)

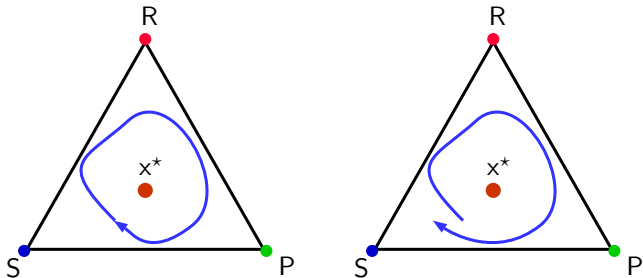
$$U^p(x^p, x^{-p}) = - \prod_{p=1}^N x^p, x^p \in [0, 1] \quad (10)$$

Game not merely monotone for $N > 2$, NE $x^* = \mathbf{0}$ is a mere VSS.

Long-run strategy generated by MD corresponds to NE in some games

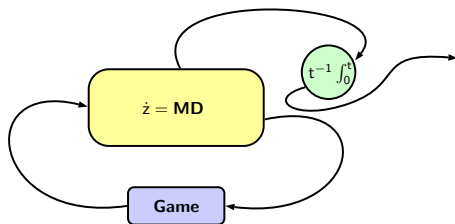
But...

MD fails to converge to NE for a wide range of games, even simple ones



Fundamental limitation in applications

Overcoming non-convergence of MD: time-averaging



Mirror descent with time-averaging (MDA) (e.g., Mertikopoulos & Sandholm, '16, Hofbauer, Sorin & Viossat '09),

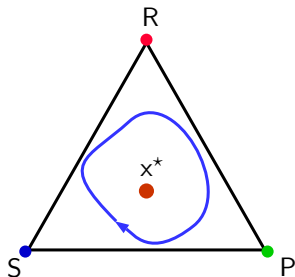
$$\dot{z} = U(x), x = C_\epsilon(z), x_{\text{avg}}(t) = t^{-1} \int_0^t x(\tau) d\tau \quad (\text{MDA})$$

- Converges exactly in ZS game with interior NE, which could be a mere VSS.
- Requires all to use time-averaging \implies not “game-realistic”.
- Not robust to slight game parameter perturbation.

Alternative: Design via passivity principles

Observation: MD fails in games with specific types of properties

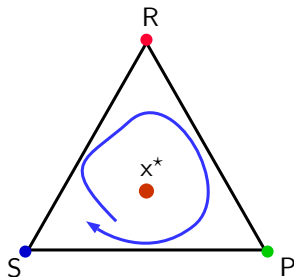
$$\begin{bmatrix} 0 & -l & w \\ w & 0 & -l \\ -l & w & 0 \end{bmatrix} \quad l = w$$



This is an example of a **merely monotone** game:

$$\underbrace{(U(x) - U(x'))^\top (x - x')}_{\text{monotonicity product}} \leq 0, \forall x, x' \in \Omega. \quad (2)$$

$$\begin{bmatrix} 0 & -l & w \\ w & 0 & -l \\ -l & w & 0 \end{bmatrix} \quad l > w$$

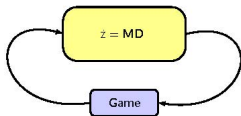


This is an example of a μ -**weakly monotone** game:

$$(U(x) - U(x'))^\top (x - x') \leq \mu \|x - x'\|_2^2, \forall x, x' \in \Omega. \quad (3)$$

Higher $\mu \geq 0$, “harder” the game, “worse” the behavior.

Observation: (weak) monotonicity of game = (lack of) energy dissipation or passivity/dissipativity of feedback subsystem U .



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3.1 Equilibrium Independent Dissipativity (EID)

Consider the system

$$\frac{d}{dt}x(t) = f(x(t), u(t)) \quad (3.3)$$

$$y(t) = h(x(t), u(t)) \quad (3.4)$$

where $x(t) \in \mathbb{R}^n$, $u(t) \in \mathbb{R}^m$, $y(t) \in \mathbb{R}^p$, and suppose there exists a set $\mathcal{X} \subset \mathbb{R}^n$ where, for every $\bar{x} \in \mathcal{X}$, there is a unique $\bar{u} \in \mathbb{R}^m$ satisfying $f(\bar{x}, \bar{u}) = 0$. Thus \bar{u} and $\bar{y} \triangleq h(\bar{x}, \bar{u})$ are implicit functions of \bar{x} .

Definition 3.1 We say that the system above is equilibrium independent dissipative (EID) with supply rate $s(\cdot, \cdot)$ if there exists a continuously differentiable storage function $V : \mathbb{R}^n \times \mathcal{X} \mapsto \mathbb{R}$ satisfying, $\forall (x, \bar{x}, u) \in \mathbb{R}^n \times \mathcal{X} \times \mathbb{R}^m$

$$V(x, \bar{x}) \geq 0, \quad V(\bar{x}, \bar{x}) = 0, \quad \nabla_x V(x, \bar{x})^T f(x, u) \leq \underbrace{s(u - \bar{u}, y - \bar{y})}. \quad (3.5)$$

From Arcak, Meissen, Packard (2016)

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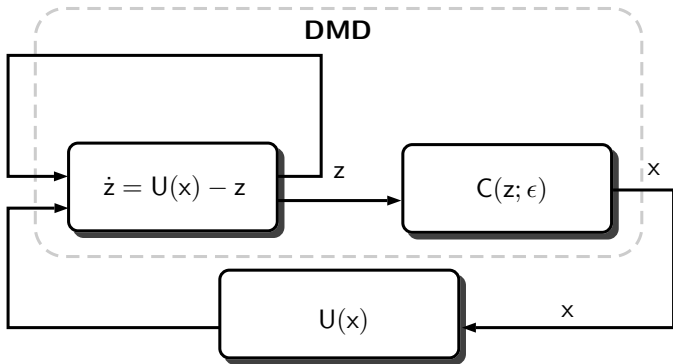
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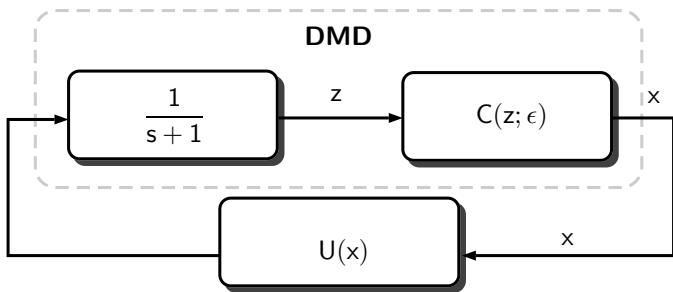
$$V(x, \bar{x}) \geq 0, \quad V(\bar{x}, \bar{x}) = 0, \quad \nabla_x V(x, \bar{x})^T f(x, u) \leq s(u - \bar{u}, y - \bar{y}). \quad (3.5)$$

From Arcak, Meissen, Packard (2016)

Idea: add **passivity/dissipativity** (more energy dissipation) to the entire system.



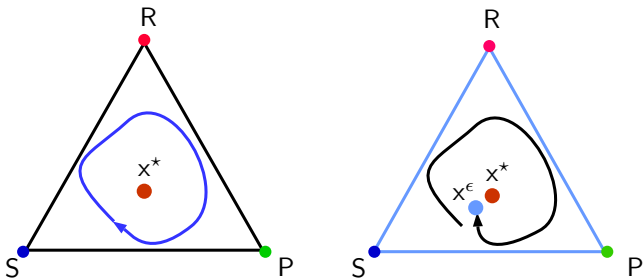
- Let players re-evaluate internal variable z during play.
- Call this **discounted mirror descent** (DMD).



DMD is “output-strictly” passive.

There exists a **storage function** V associated with DMD s.t.

$$\dot{V} \leq (x - \bar{x})^\top (U(x) - U(\bar{x})) - \epsilon \|x - \bar{x}\|_2^2$$



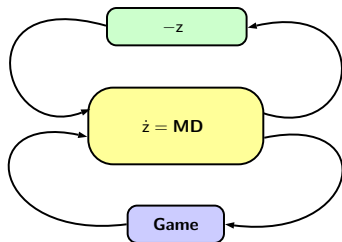
Theorem^{1,2}: \mathcal{G} merely monotone matrix game, $x(t) \rightarrow x^\epsilon$ (x^ϵ logit equilibrium).

$$x^\epsilon \rightarrow x^* \text{ as } \epsilon \rightarrow 0$$

¹B. Gao, L. Pavel, "On Passivity and Reinforcement Learning in Finite Games", in 57th IEEE CDC, 2018

²B. Gao, L. Pavel, "On Passivity, Reinforcement Learning and Higher-Order Learning in Multi-Agent Finite Games", in IEEE TAC, 2021

Overcoming non-convergence: discounting



Discounted mirror descent (DMD) (Coucheney et al. '15, Gao & Pavel '21) ,

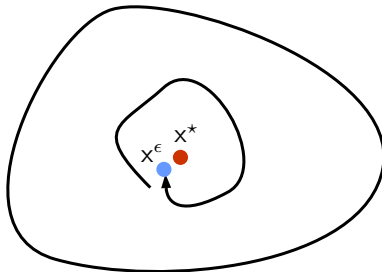
$$\dot{z} = U(x) - z, x = C_\epsilon(z) \quad \text{DMD}$$

- Converges in merely monotone game, hence to a mere VSS.
- Does not converge exactly (in general).

Reason: when dynamics settle, feedback term $-z$ does not vanish, perturbing the solution away from x^* :

$$\dot{z} = \text{DMD} = \text{MD} + \text{non-vanishing feedback}$$

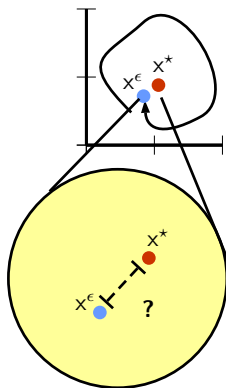
Higher-order learning MD



$x(t) \rightarrow x^*$ inexact in general

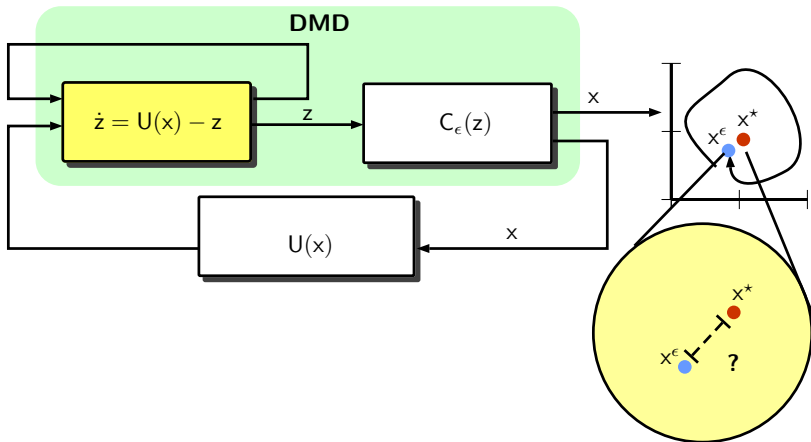
Many games lack monotonicity

Higher-Order learning MD

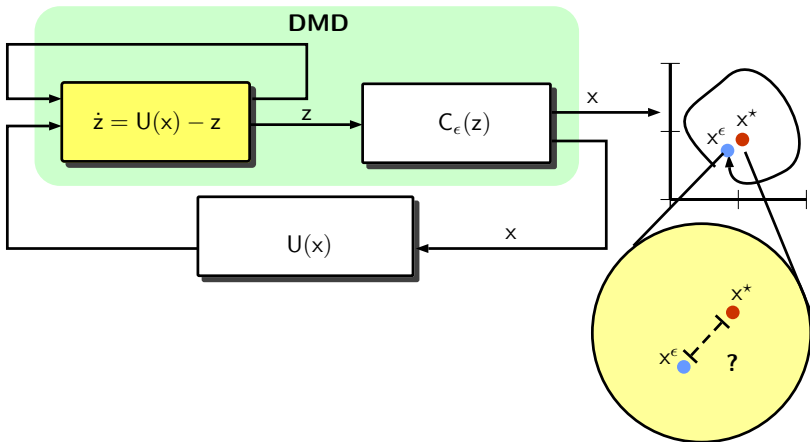


Can $x(t) \rightarrow x^*$ exactly (irrespective of ϵ)?

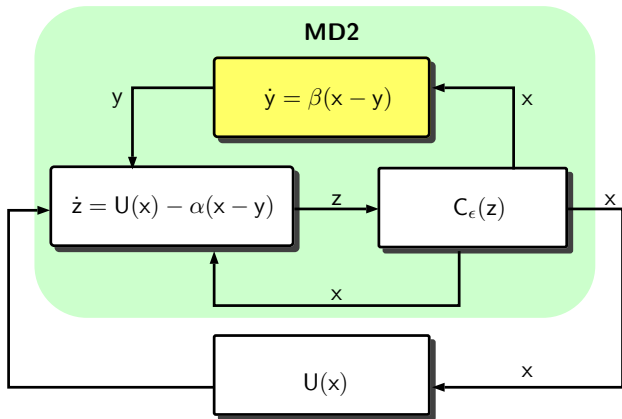
Can we do it without monotonicity?



Source of inexactness: feedback term $-z$ is **non-vanishing** at rest, perturbing the solution away from x^* .

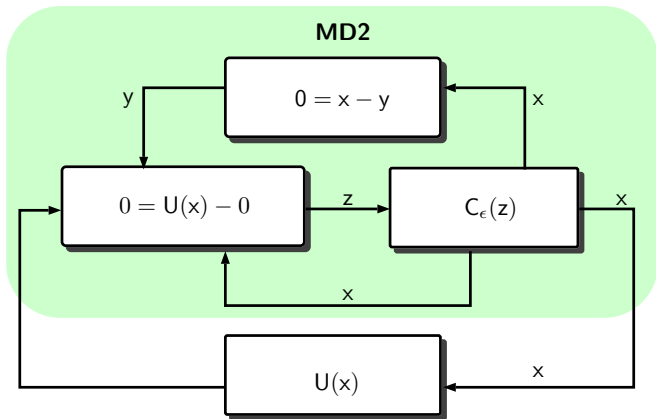


Idea: use a **new feedback** that goes away as x reaches x^* .



Introduce $\dot{y} = \beta(x - y)$, $\beta > 0$. Call it **second-order MD** (MD2).

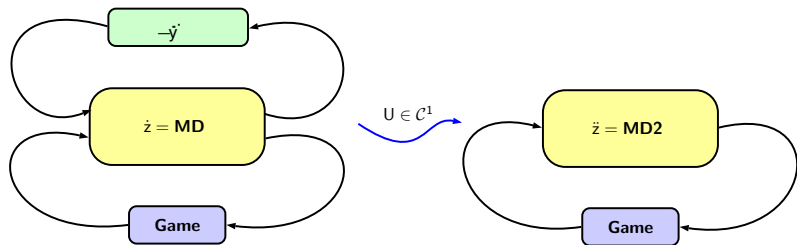
Interpretation: Do regularization but without perturbing the optimal solution.



At rest, $0 = U(x) \implies x$ an interior NE.

Why Second-Order MD? What kind of second-order?

13/



Take time derivative of \dot{z} and re-arrange,

$$\begin{aligned} \ddot{z} &= \left[\mathbf{J}_{U \circ C_\epsilon}(z) - \alpha \mathbf{J}_{C_\epsilon}(z) - \beta \mathbf{I} \right] \dot{z} + \beta U(x), \\ x &= C_\epsilon(z), \end{aligned} \quad (12)$$

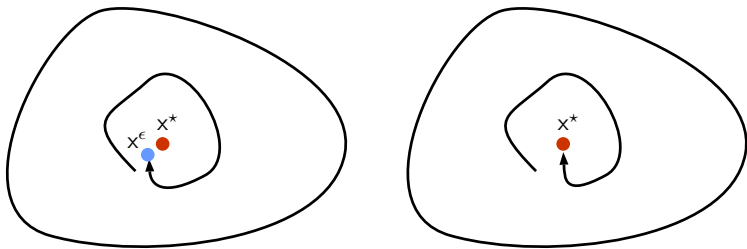
MD2 is Unlike straight second-order integration of payoffs, [Laraki & Mertikopoulos'13](#), which has the same (non-)convergence properties as MD.

MD2 \cong dual-space “heavy-ball method”, feedback modified, built on passivity-inspired principles.

Main Result: MD2 converges to interior mere VSS

Assumption. Mirror map C_ϵ^P is generated from a regularizer v^P that is \mathcal{C}^2 and either Legendre (strictly convex + boundary conditions) or strongly convex.

This assumption covers virtually all mirror maps in the literature.



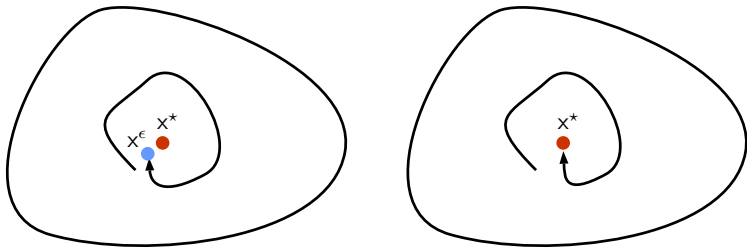
Theorem^{5,6}: $x(t) \rightarrow x^*$ **an interior mere variationally stable state (VSS)**, defined as,

$$\underbrace{U(x)^\top(x-x^*)}_{\text{variational inequality}} \leq 0, \forall x \in \Omega. \quad (4)$$

A type of NE.

⁵B. Gao and L. Pavel, "Second-order mirror descent: exact convergence beyond strictly stable equilibria in concave games," in 60th IEEE CDC, 2021.

⁶B. Gao and L. Pavel, "Second-Order Mirror Descent: Convergence in Games Beyond Averaging and Discounting," in IEEE TAC, 2024.



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Key advantage: no monotonicity of \mathcal{G} needs to be assumed.

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⁶B. Gao and L. Pavel, "Second-Order Mirror Descent: Convergence in Games Beyond Averaging and Discounting," in IEEE TAC, 2024.

MD2 Enjoys No Regret

$$\text{Regret}^P(t) = \max_{y^P \in \Omega^P} \frac{1}{t} \int_0^t \mathcal{U}^P(y^P; x^{-P}(\tau)) - \mathcal{U}^P(x(\tau)) d\tau \quad (15)$$

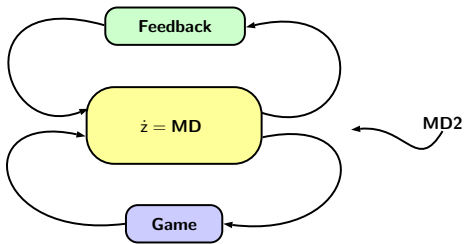
Previous result (Mertikopoulos, Papadimitriou, Piliouras, '18):
MD achieves no-regret in finite games,

$$\limsup_{t \rightarrow \infty} \text{Regret}^P(t) \leq 0, \forall p, \quad (16)$$

but cannot converge in zero-sum finite games with interior NE.

Our result: MD2 achieves no-regret in concave games with compact action sets (e.g., finite games), and converges in zero-sum games with interior mere VSS.

Second-order variant of the continuous-time mirror descent dynamics (**MD**) \Rightarrow **MD2**.



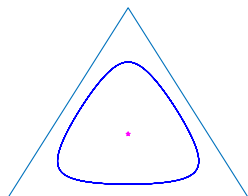
MD2 has the following benefits,

- Converges to Nash equilibrium without global game properties (e.g., (pseudo/quasi-)monotonicity)
- Converges beyond what **MD** is capable of without using additional techniques such as averaging or discounting
- Achieves exponential rate of convergence (\sim slight mod.)
- Achieves no-regret

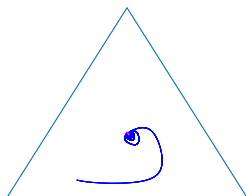
Simulation: RPS Game

$$\mathcal{A} = \begin{bmatrix} 0 & -1 & 1 \\ 1 & 0 & -1 \\ -1 & 1 & 0 \end{bmatrix} \quad \mathcal{U}^P(x^P; x^Q) = x^{P\top} \mathcal{A} x^Q$$

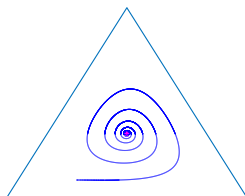
$x^* = (x^{P*}), x^{P*} = (1/3, 1/3, 1/3)$ is a mere VSS



MD



MDA

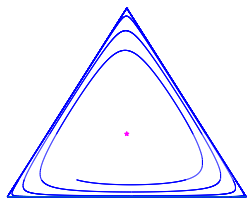


MD2

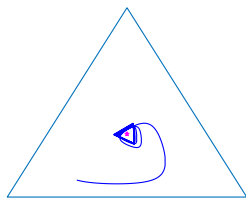
Simulation: RPS Game

$$\mathcal{A} = \begin{bmatrix} 0 & -1.2 & 1 \\ 1 & 0 & -1.2 \\ -1.2 & 1 & 0 \end{bmatrix} \quad \mathcal{U}^P(x^P; x^Q) = x^{P\top} \mathcal{A} x^Q$$

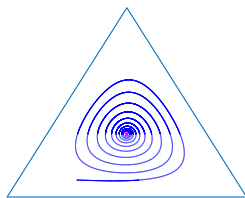
$x^* = (x^{P*}), x^{P*} = (1/3, 1/3, 1/3)$ is not a mere VSS



MD



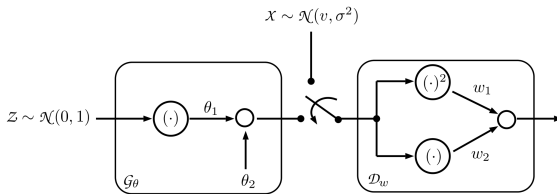
MDA



MD2

MD2 still converges (without any tuning) \Rightarrow MD2 is “robust” near the mere VSS!

Simulation: Generative Adversarial Network

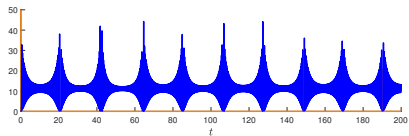


Construct $\theta = x^1, w = x^2$ s.t. $\mathcal{G}_\theta(\mathcal{Z})$ recovers mean and var of \mathcal{X} .

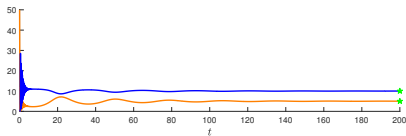
$$\mathcal{U}^1(x^1; x^2) = x_1^2(\sigma^2 + v^2 - \sum_{i=1}^2 (x_i^1)^2) + x_2^2(v - x_2^1) = -\mathcal{U}^2(x^2; x^1),$$

Game not merely monotone, but $x^* = ((\sigma, v), (0, 0))$ is mere VSS.

θ_1, θ_2 are estimates of mean and var.



MD



MD2

Parting Message

- MD2 is the systematic generalization of MD.
 - Allows to characterize discrete-time MD2 in semi-bandit and full-bandit setups.
- Future Work:
 - Systematic methods: higher-order for learning in games with coupled constraints (GNE)? (see CDC'24), continuous to discrete-time algorithms while preserving properties?
 - Other questions: distributed settings, additional constraints, discontinuity, non-convexity, even higher-orders, asynchrony, delays, time-varying, more sophisticated players...